New Constant-Weight Codes from Propagation Rules

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Abstract—This paper proposes some simple propagation rules which give rise to new binary constant-weight codes.

Index Terms—constant-weight codes, cosets, q-ary codes

I. INTRODUCTION

The ring \( \mathbb{Z}_q / q \mathbb{Z} \) is denoted \( \mathbb{Z}_q \). We endow \( \mathbb{Z}_q^n \) with the Hamming distance metric \( \Delta \): for \( u, v \in \mathbb{Z}_q^n \), \( \Delta(u, v) \) is the number of positions where \( u \) and \( v \) differ. A \((q, \text{ary})\) code of length \( n \) is a subset \( C \subseteq \mathbb{Z}_q^n \). The elements of \( C \) are called codewords, and the size of \( C \) is the number of codewords it contains. The minimum distance of a code \( C \) is \( \Delta(C) = \min_{u, v \in C, u \neq v} \Delta(u, v) \). We often denote by \((n, d)_q\)-code a q-ary code of length \( n \) and minimum distance at least \( d \).

The weight, \( wt(u) \), of \( u \in \mathbb{Z}_q^n \) is its distance from the origin, that is, \( wt(u) = \Delta(u, 0) \). For \( 0 \leq w \leq n \), the \((q, \text{ary})\) Johnson space \( J_q^n(w) \) is the set of all elements of \( \mathbb{Z}_q^n \) having weight \( w \), that is, \( J_q^n(w) = \{ u \in \mathbb{Z}_q^n : wt(u) = w \} \). A \((q, \text{ary})\) constant-weight code of length \( n \), distance \( d \), and weight \( w \), denoted \((n, d, w)_q\)-code, is a code \( C \subseteq J_q^n(w) \) such that \( \Delta(C) \geq d \).

We adopt the convention throughout this paper that if \( q \) is not specified, then we assume \( q = 2 \). Hence, for example, an \((n, d, w)_2\)-code refers to an \((n, d, w)_2\)-code, and \( J_n(w) \) refers to \( J_2^n(w) \).

Binary constant-weight codes have been extensively studied for more than four decades due to their fascinating combinatorial structures and applications [1]–[19]. Given \( n \), \( d \), and \( w \), the central problem of interest in binary constant-weight codes is in the determination of \( A(n, d, w) \), the largest possible size of an \((n, d, w)_2\)-code. Exact values of \( A(n, d, w) \) are known only for a few infinite families of parameters \( n \), \( d \), and \( w \), and in some other sporadic instances (see, for example, [3], [4]). In light of the difficulty of determining \( A(n, d, w) \) exactly, various bounds have also been developed. There are two online tables devoted to bounds on \( A(n, d, w) \): one maintained by Rain and Sloane [20] and the other by Smith and Montemanni [21]. While the former table considers codes of lengths not exceeding 63, the latter table focuses mainly on codes for lengths between 29 and 63, having small weights.

In this paper, we present simple propagation rules for binary constant-weight codes through \( q \)-ary codes. It turns out that some good binary constant-weight codes can be obtained from these propagation rules. In particular, we improve on a number of bounds in the online tables of Rain and Sloane [20], and Smith and Montemanni [21].

We remark that the table of Smith and Montemanni [21] was created because the table of Rains and Sloane [20] had not been updated for many years. For code parameters that are not covered by Smith and Montemanni [21], we have checked against recent literature, to the best of our efforts, in ascertaining that our results here do indeed improve upon existing results.

II. PROPAGATION RULES

In this section, we present some simple propagation rules for binary constant-weight codes from \( q \)-ary codes. We begin with a simple observation.

Let \( C \subseteq \mathbb{Z}_q^n \). For \( u \in \mathbb{Z}_q^n \), we denote by \( u + C \) the coset of \( C \),

\[ \{ u + c : c \in C \} \]

We also embed \( \mathbb{Z}_2 \) into \( \mathbb{Z}_q \). It is evident that \( (u + C) \cap J_q^n(w) \) is a binary constant-weight code of weight \( w \) and size \( N = |(u + C) \cap J_q^n(w)| \). Since the minimum distance \( d' \) of \( (u + C) \cap J_q^n(w) \) is at least \( d \) and \( d' \) must be even, it follows that \( d' \geq 2\lfloor (d+1)/2 \rfloor \). Thus, we have the following.

**Theorem 2.1:** Let \( 0 < w < n \). If there exists an \((n, d)_q\)-code \( C \), then there exists an \((n, 2\lfloor (d+1)/2 \rfloor, w)_2\)-code of size \( N \), where

\[ N = \max_{u \in \mathbb{Z}_q^n} |(u + C) \cap J_q^n(w)|. \]

A simple bound on the size of the constant-weight codes in Theorem 2.1 can be obtained by considering the average size of the cosets.

**Theorem 2.2:** Let \( 0 < w < n \). If there exists an \((n, d)_q\)-code of size \( M \), then

\[ A(n, 2\lfloor (d+1)/2 \rfloor, w) \geq \frac{M(n\choose w)}{q^n} \cdot \]

**Proof:** Let \( C \) be an \((n, d)_q\)-code of size \( M \). Let \( u_1, u_2, \ldots, u_q^n \) denote all the elements of \( \mathbb{Z}_q^n \), and let

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\( v_1, v_2, \ldots, v_{n^2} \) denote all the elements of \( J^n(w) \). Define
\[
\delta_{i,j} = \begin{cases} 
1, & \text{if } v_j \in u_i + C \\
0, & \text{if } v_j \notin u_i + C.
\end{cases}
\]

For each \( v_j \in J^n(w) \), there are \( M \) elements \( u_i \in \mathbb{Z}_q^n \) such that \( u_i + C \) contains \( v_j \) (to see this, note that \( v_j \in u_i + C \) if and only if \( u_i = v_j + C \) for some \( c \in C \)). Thus,
\[
\sum_{1 \leq i \leq M} \sum_{1 \leq j \leq (w^n)} \delta_{i,j} = M \binom{n}{w}.
\]
Hence, there exists at least one \( \ell \), \( 1 \leq \ell \leq q^n \), such that
\[
\sum_{1 \leq j \leq (w^n)} \delta_{i,j} \geq \frac{M \binom{n}{w}}{q^n}.
\]

The theorem now follows by noting that the size of \( (u_i + C) \cap J^n(w) \) is precisely \( \sum_{1 \leq j \leq (w^n)} \delta_{i,j} \), and we have seen above that \( (u_i + C) \cap J^n(w) \) is an \((n, 2[(d+1)/2], w)\)-code.

Next, we consider binary constant-weight codes of length \( n + 1 \) from \( q \)-ary codes of length \( n \).

**Theorem 2.3:** Let 0 < \( w < n \). Suppose there exists an \((n, d)_q\)-code \( C \) of size \( M \). Then,
(i) there exists an \((n+1, 2[(d+1)/2], w)\)-code of size \( N \), where
\[
N = \max_{u \in \mathbb{Z}_q^n} |(u + C) \cap (J^n(w-1) \cup J^n(w))|;
\]
(ii)
\[
A(n+1, 2[(d+1)/2], w) \geq \left\lceil \frac{M \binom{n}{w-1} + \binom{n}{w}}{q^n} \right\rceil.
\]

**Proof:**
(i) Let \( u \in \mathbb{Z}_q^n \) such that \(|(u + C) \cap (J^n(w-1) \cup J^n(w))|\) achieves the maximum size \( N \). It is clear that \( C' = (u + C) \cap (J^n(w-1) \cup J^n(w)) \) is an \((n, d)\)-code, where each codeword has weight either \( w-1 \) or \( w \). To each codeword \( c \in C' \), append a new coordinate which takes on value one if \( wt(c) = w-1 \) and value zero if \( wt(c) = w \). The set of resulting codewords is an \((n+1, 2[(d+1)/2], w)\)-code.
(ii) Using the same arguments as in the proof of Theorem 2.2, we get an \((n, d)\)-code of size \( M \binom{n}{w-1} + \binom{n}{w}/q^n \), in which the weight of every codeword is either \( w-1 \) or \( w \). By appending a new coordinate to every codeword as in (i) above, we get an \((n+1, 2[(d+1)/2], w)\)-code of the required size.

### III. Examples

We provide some examples where the propagation rules given by Theorems 2.2 and 2.3 lead to improved bounds on \( A(n, d, w) \).

In the tables of this section, a bold entry indicates that the size of the code constructed here is larger than any known codes of the same parameters, and a entry superscripted by an asterisk indicates that the size of the code constructed here is of the same size as the best known code of the same parameters. \( M_{\text{max}} \) denotes the lower bound on \( A(n, d, w) \) given by Theorems 2.1 or 2.3(ii), and \( M_{\text{avg}} \) denotes the lower bound on \( A(n, d, w) \) given by Theorems 2.2 or 2.3(ii). \( M_{\text{BS}} \) denotes the lower bound on \( A(n, d, w) \) in the tables of Rains and Sloane [20].

**Example 3.1:** Let \( C \) be the Goethals \((63, 7)\)-code of size \( 2^{47} \) [22] (see [23, Chapter 5] for the structure of this code).
- Theorems 2.2 and 2.3(ii) give
\[
A(63, 8, w) \geq \left\lceil \frac{63}{w} / 2^{16} \right\rceil,
\]
\[
A(64, 8, w) \geq \left\lceil \left( \frac{63}{w-1} + \frac{63}{w} \right) / 2^{16} \right\rceil.
\]
The implications of these bounds are given in Table I.

<table>
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<th>( M_{\text{max}} )</th>
<th>( M_{\text{BS}} )</th>
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- Shortening \( C \) at the last \( i \) positions, \( 1 \leq i \leq 46 \), results in a \((63 - i, 7)\)-code of size \( 2^{47-i} \). It follows from Theorem 2.2 that there exists a \((63 - i, 8, 7)\)-code of size \( \binom{63-i}{7}/2^{16} \). In particular, when \( i \in \{1, 2, 3\} \), this implies
\[
A(62, 8, 7) \geq 7505,
\]
\[
A(61, 8, 7) \geq 6657,
\]
\[
A(60, 8, 7) \geq 5894.
\]
The three lower bounds (1)–(3) improve those in [21] (the corresponding lower bounds given there are 6693, 6223, and 5770, respectively, obtained by Smith et al. [13]).

**Example 3.2:** Let \( C \) be the Preparata \((63, 5)\)-code of size \( 2^{52} \) [24] (see [23, Chapter 5] for the structure of this code). Theorems 2.2 and 2.3(ii) give
\[
A(63, 6, 8) \geq \left\lfloor \frac{63}{w} / 2^{11} \right\rfloor,
\]
\[
A(64, 6, w) \geq \left\lfloor \left( \frac{63}{w-1} + \frac{63}{w} \right) / 2^{11} \right\rfloor.
\]
We also found via computation cosets of \( C \) achieving the maximum in Theorems 2.1 and 2.3(ii). The results are given in Tables II and III.

**Example 3.3:** Let \( C \) be the \((31, 9)\)-code of size \( 2^{13} \) constructed by Grassl [25].
- We found via computation cosets of \( C \) achieving the maximum in Theorems 2.1 and 2.3(i). The results are given in Table IV.
We found, via computation, cosets of \(C\) in a \(w\)-code of size \(#C\). We found, via computation, cosets of this shortened code achieving the maximum in Theorem 2.3(i). The results are not known for these parameter sets.

Example 3.5: Let \(C\) be the (linear) \((31,13)\)-code of size \(2^7\) constructed by Grassl [25]. We found, via computation, cosets of \(C\) achieving the maximum in Theorem 2.3(i). These provide the lower bounds

\[
A(32, 14, 12) \geq 29, \\
A(32, 14, 13) \geq 42.
\]

Lower bounds on \(A(32, 14, w)\), \(w \in \{12, 13\}\), are previously not known.

Example 3.6: Let \(C_0\) be the (linear) Reed-Muller \((32,16)\)-code of size \(2^6\), and let \(C\) be the code obtained from \(C_0\) by puncturing it at the last position. Then \(C\) is a \((31,15)\)-code of size \(2^6\). We found, via computation, cosets of \(C\) achieving the maximum in Theorems 2.1 and 2.3(i). These provide the lower bounds

\[
A(n, 16, 13) \geq 16, \\
A(n, 16, 14) \geq 21, \\
A(n, 16, 15) \geq 31,
\]

for \(n \in \{31,32\}\). Lower bounds on \(A(n, 16, w)\), \(w \in \{12, 13\}\), are previously not known for these parameters.

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